

On the Informed Principal Model with Common Values

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Motivation (I)

- Conventional mechanism design theory assumes that the party that designs a mechanism (i.e., the principal) has no private (payoff-relevant) information
- However, in many occasions the principal may have private information. For instance:
 - ▶ **Procurement:** The government may have superior information about the cost of a project than potential constructors
 - ▶ **Vertical contracting:** An upstream manufacturer may have more detailed information at the time of contracting about market characteristics, e.g., the market demand, than a downstream retailer
 - ▶ **Informed seller:** A seller of an indivisible object may have superior information about the quality of the object than potential buyers

Motivation (II)

- In all these occasions, the principal has private information that affects the value of trade
- When offering the mechanism the principal may reveal part of this information and therefore the design of the mechanism itself becomes subtle
- Question: What mechanism will the principal select?
- Myerson (1983) took an axiomatic approach to characterise a reasonable solution
- Maskin and Tirole (1990, 1992) (MT) took a non-cooperative approach to characterise mechanisms that can result as equilibria in a three stage game (mechanism proposal/acceptance-rejection/mechanism execution)

Contribution

- I identify some fundamental properties of the Rothschild-Stiglitz-Wilson allocation, i.e., the undominated allocation within the set of incentive compatible and individually rational for the agent type by type allocations
- Based on these properties, I construct a more robust, and perhaps simpler, proof of Theorem 1 (the main result) of Maskin and Tirole (1992)
- I make a distinction between simple mechanisms (in which the agent makes no announcement, e.g., DRMs) and general mechanisms
- I provide a simple example to highlight why such distinction is important
- I provide a more general condition than the no-tangency condition provided in Maskin and Tirole (1992) that allows for the complete characterisation of the set of equilibrium allocations

The Model with Unilateral Private Information

- Two players: a principal (P) and an agent (A)
- Type of principal $i = 1, \dots, n$, $n \geq 2$, is her private information
- The set of actions is $X \subset \mathbb{R}^K$, $K \geq 1$ (compact)
- Prior beliefs: $\Pi = (\Pi^i)_i$, where $\Pi^i > 0$ for every i
- Payoffs: $V^i(x)$ and $U^i(x)$ (continuous) for the principal and the agent respectively
- The two players wish to select a (potentially random) contractible action from $\mathcal{M} = \Delta(X)$
- Let $V^i(\mu) = \int V^i(x)d\mu$ and $U^i(\mu) = \int U^i(x)d\mu$

Allocations - IC, Dominance, IE

- Allocation: $\mu^i = (\mu^i)_i$

Definition

An allocation μ^i is incentive compatible (IC) if $U^i(\mu^i) \geq U^i(\mu^j)$ for every i, j

Definition

An IC allocation μ^i dominates an IC allocation $\bar{\mu}^i$ if $V^i(\mu^i) \geq V^i(\bar{\mu}^i)$ for every i with the inequality being strict for at least one i

Definition

An allocation $\bar{\mu}^i$ is interim efficient (IE) relative to beliefs $\bar{\Pi}^i$ (not necessarily the prior beliefs) if (i) it is IC, and (ii) there exists no allocation $\mu^i \neq \bar{\mu}^i$ that is IC, satisfies $\sum \bar{\Pi}^i U^i(\mu^i) \geq \sum \bar{\Pi}^i U^i(\bar{\mu}^i)$, and dominates $\bar{\mu}^i$.*

*An allocation is weakly interim efficient (WIE) (or IE type by type) if we substitute $\sum \bar{\Pi}^i U^i(\mu^i) \geq \sum \bar{\Pi}^i U^i(\bar{\mu}^i)$ with $U^i(\mu^i) \geq U^i(\bar{\mu}^i)$ for every i .

Reservation Allocation - IR

- *Reservation Allocation*: μ_0^i ; it can be regarded as either an outside option or a prior contract that binds the two players and they wish to renegotiate.

Definition

An IC allocation μ^i is individually rational (IR) relative to beliefs $\bar{\Pi}^i$ (not necessarily the prior beliefs) if $\sum \bar{\Pi}^i U^i(\mu^i) \geq \sum \bar{\Pi}^i U^i(\mu_0^i)$.*

* An allocation μ^i is IR type by type if we substitute

$$\sum \bar{\Pi}^i U^i(\mu^i) \geq \sum \bar{\Pi}^i U^i(\mu_0^i) \text{ with } U^i(\mu^i) \geq U^i(\mu_0^i) \text{ for every } i.$$

The Rothschild-Stiglitz-Wilson (RSW) Allocation

Definition

An allocation μ^i is an RSW allocation (relative to the reservation allocation μ_0^i) if (i) it is IC and IR type by type, and (ii) there exists no allocation $\mu^i \neq \bar{\mu}^i$ that is IC, IR type by type, and dominates μ^i .

The Rothschild-Stiglitz-Wilson (RSW) Allocation

Lemma

Every RSW allocation is payoff equivalent for the principal.

Proof.

Suppose that $\hat{\mu}_1(\mu_0)$ and $\hat{\mu}_2(\mu_0)$ are RSW allocations, where $\hat{\mu}_1(\mu_0) \neq \hat{\mu}_2(\mu_0)$ and $V^i(\hat{\mu}_1^i(\mu_0)) \neq V^i(\hat{\mu}_2^i(\mu_0))$ for some i . Let $I_1 = \{i : V^i(\hat{\mu}_1^i(\mu_0)) \geq V^i(\hat{\mu}_2^i(\mu_0))\}$ and $I_2 = \{i : V^i(\hat{\mu}_1^i(\mu_0)) < V^i(\hat{\mu}_2^i(\mu_0))\}$. Because $\hat{\mu}_1(\mu_0)$ and $\hat{\mu}_2(\mu_0)$ are IC and IR type by type, allocation $\tilde{\mu}$, which maps each type from I_1 to her action in allocation $\hat{\mu}_1(\mu_0)$ and each type from I_2 to her action in allocation $\hat{\mu}_2(\mu_0)$, is also IC and IR type by type. Allocation $\tilde{\mu}$ dominates both $\hat{\mu}_1(\mu_0)$ and $\hat{\mu}_2(\mu_0)$, which contradicts the definition of an RSW allocation. \square

Mechanisms

- *Mechanism*: $m = (S, g)$, $S = S_P \times S_A$ (finite), $g : S \rightarrow \mathcal{M}$

Definition

For given beliefs $\bar{\Pi}$, a Bayesian Nash equilibrium in mechanism m consists of a profile of strategies, one for each player, such that, conditional on the strategy of the other player, no player has a unilateral profitable deviation.

- Every equilibrium in mechanism m under beliefs $\bar{\Pi}$ is associated with an ex post (expected) equilibrium payoff profile $(\bar{V}(m, \bar{\Pi}), \bar{U}(m, \bar{\Pi}))$, where $\bar{V}(m, \bar{\Pi}) = (\bar{V}^i(m, \bar{\Pi}))_i$ and $\bar{U}(m, \bar{\Pi}) = (\bar{U}^i(m, \bar{\Pi}))_i$
- Simpler class of mechanisms is the class of *direct revelation mechanisms* (DRMs), in which the principal simply announces a type (not necessarily the true type), and the agent makes no announcement.

Properties of the RSW Allocation - General Mechanisms

Proposition

Suppose that $\hat{\mu}(\mu_0)$ (i.e., the RSW allocation) is IE relative to beliefs $\hat{\Pi}$ (not necessarily the prior beliefs), where $\hat{\Pi}^i > 0$ for every i ; then, for every mechanism $m \neq \hat{\mu}(\mu_0)$ and subset of types $I \subseteq \{1, \dots, n\}$, there exist beliefs $\bar{\Pi}$ such that in every equilibrium of m under $\bar{\Pi}$ with an associated equilibrium payoff profile $(\bar{V}(m, \bar{\Pi}), \bar{U}(m, \bar{\Pi}))$, either

$$\bar{V}^i(m, \bar{\Pi}) \leq V^i(\hat{\mu}^i(\mu_0)) \text{ for every } i \in I \quad (1)$$

or

$$\sum \bar{\Pi}^i \bar{U}^i(m, \bar{\Pi}) < \sum \bar{\Pi}^i U^i(\mu_0^i) \quad (2)$$

Proof of Proposition 1

Suppose that there exists $m \neq \hat{\mu}^i(\mu_0)$ and $I \subseteq \{1, \dots, n\}$ such that for every $\bar{\Pi}^i$, there exists an equilibrium with an associated equilibrium payoff profile $(\bar{V}^i(m, \bar{\Pi}^i), \bar{U}^i(m, \bar{\Pi}^i))$, such that: (i) $\bar{V}^i(m, \bar{\Pi}^i) > V^i(\hat{\mu}^i(\mu_0))$ for every $i \in I$, and $\bar{V}^i(m, \bar{\Pi}^i) \leq V^i(\hat{\mu}^i(\mu_0))$ for every $i \notin I$, and, (ii) $\sum \bar{\Pi}^i \bar{U}^i(m, \bar{\Pi}^i) \geq \sum \bar{\Pi}^i U^i(\mu_0^i)$

Consider Π_i , where $\Pi_i^j = \hat{\Pi}^j / \sum_{j \in I} \hat{\Pi}^j$ for every $i \in I$, and $\Pi_i^j = 0$ for every $i \notin I$.

Construct $\tilde{\mu}^i$, where $V^i(\tilde{\mu}^i) = \bar{V}^i(m, \Pi_i)$ for every $i \in I$ and $\tilde{\mu}^i = \hat{\mu}^i(\mu_0)$ for every $i \notin I$. $\tilde{\mu}^i$ is IC because of (i) and $\hat{\mu}^i(\mu_0)$ is IC by definition. Moreover,

$$\sum_{i \in I} \hat{\Pi}^i U^i(\tilde{\mu}^i) = \sum_{i \in I} \hat{\Pi}^i \bar{U}^i(m, \Pi_i) + \sum_{i \notin I} \hat{\Pi}^i U^i(\hat{\mu}^i(\mu_0)) \geq \sum \hat{\Pi}^i U^i(\mu_0^i) \quad (3)$$

because $\sum_{i \in I} \frac{\hat{\Pi}^i}{\sum_{j \in I} \hat{\Pi}^j} \bar{U}^i(m, \Pi_i) \geq \sum_{i \in I} \frac{\hat{\Pi}^i}{\sum_{j \in I} \hat{\Pi}^j} U^i(\mu_0^i)$ from (ii) above and due to the fact that the RSW allocation is IR type by type. $\tilde{\mu}^i$ dominates $\hat{\mu}^i(\mu_0)$ and is IR relative to beliefs $\hat{\Pi}^i$; a contradiction.

Properties of the RSW Allocation - Simple Mechanisms

Proposition

Suppose that the principal is restricted to offering only DRMs; then, for every IC allocation $\mu^i \neq \hat{\mu}^i(\mu_0^i)$ (where $\hat{\mu}^i(\mu_0^i)$ is the RSW allocation) and subset of types $I \subseteq \{1, \dots, n\}$, there exist beliefs $\bar{\pi}^i$ such that either

$$V^i(\mu^i) \leq V^i(\hat{\mu}^i(\mu_0^i)) \quad \text{for every } i \in I \quad (4)$$

or

$$\sum \bar{\pi}^i U^i(\mu^i) < \sum \bar{\pi}^i U^i(\mu_0^i) \quad (5)$$

Proof of Proposition 2

Suppose that there exists an IC allocation $\mu^i \neq \hat{\mu}^i(\mu_0^i)$ and $I \subseteq \{1, \dots, n\}$ such that for every $\bar{\Pi}^i$: (i) $V^i(\mu^i) > V^i(\hat{\mu}^i(\mu_0^i))$ for every $i \in I$ and $V^i(\mu^i) \leq V^i(\hat{\mu}^i(\mu_0^i))$ for every $i \notin I$, and (ii) $\sum \bar{\Pi}^i U^i(\mu^i) \geq \sum \bar{\Pi}^i U^i(\mu_0^i)$. Consider $\tilde{\mu}^i$, where

$$\tilde{\mu}^i = \begin{cases} \mu^i, & \text{if } i \in I \\ \hat{\mu}^i(\mu_0^i), & \text{otherwise} \end{cases}$$

This allocation is IC because μ^i and $\hat{\mu}^i(\mu_0^i)$ are IC and (i) above. Moreover, $\tilde{\mu}^i$ is IR type by type because $\hat{\mu}^i(\mu_0^i)$ and μ^i are IR type by type (by the definition of the RSW allocation and (ii) above). Therefore, μ^i dominates $\hat{\mu}^i(\mu_0^i)$ and is IR type by type, which contradicts $\hat{\mu}^i(\mu_0^i)$ being an RSW allocation.

Example 1

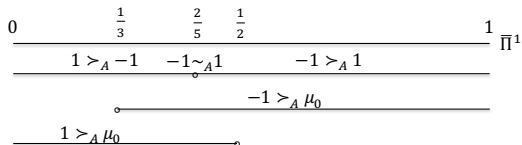
- $n = 2$, $\Pi^1 = \Pi^2 = 1/2$ and $\mathcal{M} = \{-1, 0, 1\}$

μ	$U^1(\mu)$	$U^2(\mu)$	$V^1(\mu)$	$V^2(\mu)$
-1	2	-1	2	2
0	0	0	1	1
1	-1	1	3	3
μ_0	0	0	0	0

Table: Payoffs

- IC allocations: $\{(-1, -1), (0, 0), (1, 1)\}$
- RSW allocation: $\hat{\mu}(\mu_0) = (0, 0)$

Example 1 cont'd



- RSW is not IE relative to any non-degenerate beliefs
- Consider mechanism $m^d = (S^d, g^d)$, where $S_P^d = \emptyset$, $S_A^d = \{-1, 1\}$ and $g^d(s) = s$ for every $s \in S_A^d$
- For DRM $(-1, -1)$ and $\bar{\Pi}^1 < 1/3$, $\sum \bar{\Pi}^i U^i(-1) < 0$; for DRM $(1, 1)$ and $\bar{\Pi}^1 > 1/2$, $\sum \bar{\Pi}^i U^i(1) < 0$

The Extensive-Form Game

- As in MT let the two players play the following game
 - 1 The principal proposes a mechanism
 - 2 The agent accepts or rejects
 - 3 If the agent rejects, the reservation action is in effect. If he accepts, the two players play the mechanism proposed by the principal
- Strategies of players
- Update of beliefs
- *Perfect Bayesian Equilibrium*
- **Inscrutability Principle (Myerson 1983)**: There is no loss of generality in concentrating on equilibria in which all types offer the same mechanism

Equilibrium Allocations

Suppose that X is convex and a type i indifference curve is nowhere tangent to a type j indifference curve

Theorem (MT92)

Suppose that $\hat{\mu}^i(\mu_0^i)$ (i.e., the RSW allocation) is IE relative to beliefs $\hat{\Pi}^i$ (not necessarily the prior beliefs), where $\hat{\Pi}^i > 0$ for every i ; then, an allocation $\bar{\mu}^i$ is an equilibrium allocation of the three-stage game if and only if it is IC and satisfies

$$V^i(\bar{\mu}^i) \geq V^i(\hat{\mu}^i(\mu_0^i)) \quad \forall i \quad (6)$$

$$\sum \Pi^i U^i(\bar{\mu}^i) \geq \sum \Pi^i U^i(\mu_0^i) \quad (7)$$

Sketch of Proof of Theorem 1

* For the “if” part:

- The sets of mechanisms can be partitioned in two subsets: Set A includes those mechanisms that are IR for the agent relative to all possible beliefs; Set A^c includes every other feasible mechanism
- Suppose that the principal offers $\bar{\mu}$; an IC allocation that satisfies (8) and (7)
- To construct an equilibrium one needs to assign beliefs to every feasible mechanism and specify sequentially rational strategies

Sketch of Proof of Theorem 1 cont'd

- For every mechanism in A , Proposition 1 assures that there exist beliefs such that all types are worse off relative to the RSW allocation
- For every mechanism in A^c , assign beliefs such that the mechanism is not IR for the agent
- Construct sequentially rational strategies for the principal and the agent given these beliefs
- Mission accomplished!

* For the “only if” part key is the assumption that the indifference curves of different types are nowhere tangent.

Robustness

- Maskin and Tirole (1992) show that for every mechanism different from the on-the-equilibrium path mechanism there exist beliefs and **an** equilibrium (i.e., continuation of the game) such that every type is worse off
- The proof proposed here shows that for every mechanism different from the on-the-equilibrium path mechanism there exist beliefs such that in **every** equilibrium (i.e., continuation of the game) every type is worse off

A More General Sufficient Condition for the “Only If”

- The “nowhere no-tangency” condition might be too strong
- Is there a milder condition that allows for the complete characterisation of the set of equilibrium allocations?
- Two further definitions

Definition

An allocation μ^i is strictly incentive compatible (SIC) if $V^i(\mu^i) > V^i(\mu^j)$ for every i, j .

Definition

An IC allocation μ^i is strictly individually rational (SIR) type by type if $U^i(\mu^i) > U^i(\mu_0^i)$.

A More General Sufficient Condition for the “Only If”

Proposition

Suppose that there exists a sequence of SIC and SIR type by type allocations $\{\mu_p^i\}_{p=1}^{\infty}$, i.e., $V^i(\mu_p^i) > V^i(\mu_p^j)$ for every i, j, p and $U^i(\mu_p^i) > U^i(\mu_0^i)$ for every i, p , such that $\{V^i(\mu_p^i)\}_{p=1}^{\infty}$ converges to $V^i(\hat{\mu}^i(\mu_0^i))$ for every i ; then every equilibrium allocation of the three-stage game $\bar{\mu}^i$ is such that

$$V^i(\bar{\mu}^i) \geq V^i(\hat{\mu}^i(\mu_0^i)) \quad \forall i \quad (8)$$

Proof of Proposition 3

Consider an IC allocation $\bar{\mu}^j$ in which for some type j , $V^j(\bar{\mu}^j) < V^j(\hat{\mu}^j(\mu_0^j))$. Let $V^j(\hat{\mu}^j(\mu_0^j)) - V^j(\bar{\mu}^j) = \delta > 0$. The following lemma facilitates the proof.

Lemma

There exists p_δ such that $V^j(\mu_p^j) \geq V^j(\bar{\mu}^j)$ for every $p \geq p_\delta$.

Proof of Lemma. Because there exists a sequence $\{\mu_p^i\}_{p=1}^\infty$ such that $\{V^i(\mu_p^i)\}_{p=1}^\infty$ converges to $V^i(\hat{\mu}^i(\mu_0^i))$ for every i , there exists p_δ such that $|V^j(\mu_p^j) - V^j(\hat{\mu}^j(\mu_0^j))| < \delta$ for every $p \geq p_\delta$. Suppose that $V^j(\mu_p^j) > V^j(\hat{\mu}^j(\mu_0^j))$ for some $p \geq p_\delta$. Consider $\tilde{\mu}^j$, where $\tilde{\mu}^j = \mu_p^j$ for $i = j$ and $\tilde{\mu}^i = \hat{\mu}^i(\mu_0^i)$ for $i \neq j$. Allocation $\tilde{\mu}^j$ is IC and IR type by type which contradicts the definition of the RSW allocation. Therefore, $V^j(\mu_p^j) \leq V^j(\hat{\mu}^j(\mu_0^j))$ for every $p \geq p_\delta$. Then, $V^j(\hat{\mu}^j(\mu_0^j)) - V^j(\mu_p^j) < \delta$ and hence $V^j(\hat{\mu}^j(\mu_0^j)) - V^j(\mu_p^j) < V^j(\hat{\mu}^j(\mu_0^j)) - V^j(\bar{\mu}^j)$, which is equivalent to $V^j(\mu_p^j) > V^j(\bar{\mu}^j)$.

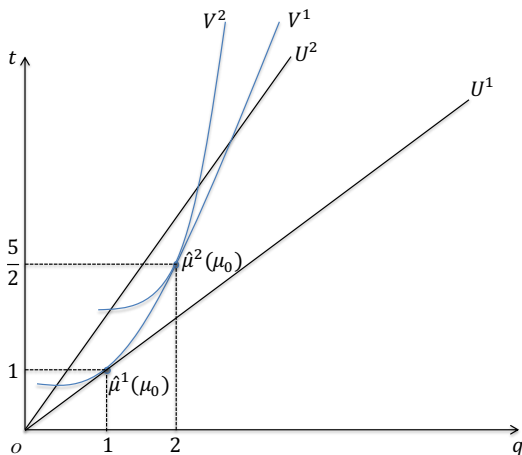
Proof of Proposition 3 cont'd

Consider mechanism $\mu_{p'_\delta}$, where $p'_\delta \geq p_\delta$. Because this mechanism is SIC and SIR type by type, it provides the agent with a payoff strictly greater than the payoff he can obtain in the reservation allocation regardless of his beliefs. Therefore, if this mechanism is proposed by type j , it should be accepted by the agent; otherwise the equilibrium fails to be sequentially rational. Type j can achieve a higher payoff by proposing mechanism $\mu_{p'_\delta}$ than by proposing $\bar{\mu}$, which means that allocation $\bar{\mu}$ cannot constitute an equilibrium allocation.

Example 2

- $n=2, K = 2$
- $V^1(t, q) = t - q^2/2, V^2(t, q) = t - 2q(1 + q)/5$
- $MRS_{t,q}^1 = -\frac{v_q^1}{V_t^1} = q, MRS_{t,q}^2 = -\frac{v_q^2}{V_t^2} = 2(1 + 2q)/5$
- Indifference curves are tangent at $q = 2, MRS_{t,q}^1 > MRS_{t,q}^2$ if $q > 2$ and $MRS_{t,q}^1 < MRS_{t,q}^2$ if $q < 2$
- RSW allocation is $((1, 1), (2, 5/2))$

Example 2 cont'd



Example 2 cont'd

- Consider $\mu_p^1 = (1 - 1/p, 1)$ and $\mu_p^2 = (5/2 - 2/p, 2)$

$$V^1(t_p^1, q_p^1) = 1/2 - 1/p > 1/2 - 2/p = V^1(t_p^2, q_p^2) \text{ for every } p$$

$$V^2(t_p^2, q_p^2) = -51/30 - 2/p > 1/5 - 1/p = V^2(t_p^1, q_p^1) \text{ for every } p$$

$$U^1(t_p^1, q_p^1) = 1/p > 0, U^2(t_p^2, q_p^2) = 3/2 + 2/p > 0 \text{ for every } p$$

- Hence μ_p^i is SIC and SIR for every p and

$$\{V^i(t_p^i, q_p^i)\}_p \rightarrow V^i(\hat{\mu}^i(\mu_0)) \text{ for every } i$$

Take-aways

- In general environments, restriction to DRMs restricts the set of profitable deviations and therefore allows one to establish the existence of equilibrium
- However, this is with loss of generality
- The revelation principle holds (i.e., restriction to DRMs is without loss of generality) iff the RSW allocation is IE relative to some non-degenerate beliefs

Future Work

- Better characterisation of environments that satisfy the sufficient condition that allows for the complete characterisation of the set of equilibrium allocations
- Better characterisation of environments in which the RSW allocation is IE relative to some non-degenerate beliefs
- Environments with bilateral private information
- Applications